Day 22. Quadrature, cont'd

culminating in adaptive quadrature with Gauss-Kronrod rules

Recall that we wish to accurately approximate \int_{a}^{b}

in the common situation where we do not have access to an antiderivative F, F' = f, hence cannot used the Fund. Thm. of Calculus.

We will use a "quadrature rule"

$$Q_{a,j}(f) = q\left(\left\{f(x_j)\right\}\right) = \sum_{j=0}^{m} \omega_j f(x_j) = (b-a) \sum_{j=0}^{m} \kappa_j f(x_j)$$

A quadrature rule is said to be "of polynomial degree m" if

$$Q_{(a,b)}(p) = \int_{a}^{b} P \quad \forall p \in \mathbb{P}$$

If the m+1 $\{x, y, y\}_{y,z,0}$ are already chosen and distinct, then the requirement of Q being of polynomial degree m consitutes m+1 conditions on the m+1 weights $\{x, y\}_{y,z,0}$ and we can determine those weights in this way.

d we can determine those weights in this way.

$$Q(1) \stackrel{\text{require}}{=} \int_{a}^{b} dx \qquad (b-a) \stackrel{\text{S}}{>} | \alpha_{j} = b-a$$

$$Q(x) = \int_{a}^{b} dx \qquad (b-a) \stackrel{\text{S}}{>} | \alpha_{j} = \frac{b^{2}-a^{2}}{2}$$

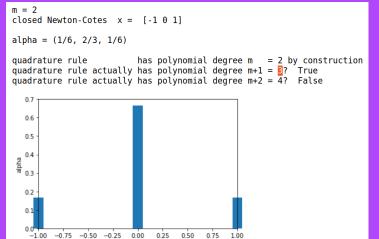
$$Q(x^{2}) = \int_{a}^{b} x^{2} dx \qquad (b-a) \stackrel{\text{S}}{>} | \alpha_{j} = \frac{b^{3}-a^{3}}{3}$$

$$Q(x^{m}) = \int_{a}^{b} x^{m} dx \qquad (b-a) \stackrel{\text{S}}{>} | \alpha_{j} = \frac{b^{m+1}-a^{m+1}}{m+1}$$

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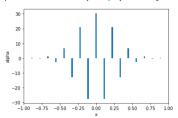
m+1 linear equations in m+1 variables $\{x_j\}$.
Since the matrix is non-singular for distinct $\{x_j\}$, with m+1 sample points we can achieve polynomial degree at least m.

Let's look at some examples ...

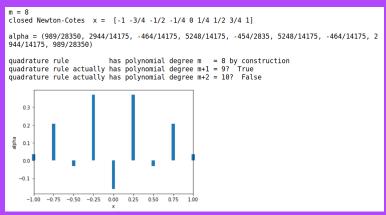


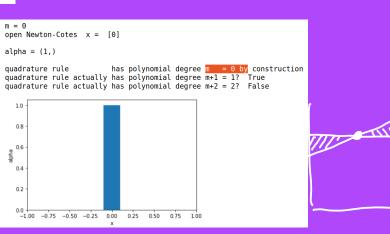
alpha = (203732352169/15209113920000, 6162434073/50697046400, -214182958293/1013940928000, 1 61769065751/158428270000, . 716535961191/63371308000, 3085731055533/45265220000, . 23648622603 3/18106688000, 6646578843333/31685654000, -69854658519033/2534852320000, 153183485240000, . 169854658519033/2534852320000, 664657884333/31685654000, -236486226033/18106088000, 308573105533/45265220000, . 176535961191/63371308000, 161769065751/158428270000, -21418295829 3/1013940928000, 16162434073/50697046400, 203732352169/15209113920000

quadrature rule actually has polynomial degree m = 18 by construction quadrature rule actually has polynomial degree m+1 = 19? True quadrature rule actually has polynomial degree m+2 = 20? False



Large alternating positive and negative weights can lead to catastrophic loss of significance due to subtraction of near-equals.





```
m = 2 open Newton-Cotes x = [-1/2 \ 0 \ 1/2] alpha = (2/3, -1/3, 2/3) quadrature rule has polynomial degree m = 2 by construction quadrature rule actually has polynomial degree m+1 = 3? True quadrature rule actually has polynomial degree m+2 = 4? False 0.6 - \frac{0.6}{0.4} - \frac{0.2}{\frac{20}{6}} = 0.2 - \frac{0.2}{\frac{20}{6}} = 0.2 - \frac{0.2}{0.00} = 0.25 - 0.00 = 0.25 = 0.50 = 0.75 = 100
```

We see the results with uniformly spaced nodes are not terribly satisfactory, especially in regard to having negative weights.

But who said the nodes should be uniformly spaced?

How about considering node locations to be another set of degrees of freedom? Then we'd have 2m+2 variables: perhaps we could then achieve polynomial degree 2m+1 / ?

Turns out this IS possible, and the corresponding optimal rules are called Gaussian quadrature.

Example: the midpoint rule, aka 1-point open Newton-Cotes, is also the 1-point Gaussian rule:

$$m=0$$
, $x_0=\frac{a+b}{2}$, $x_0=1$ has polynomial degree $2.0+1=1$. = $m+1$.

And the 4-point Gaussian rule has degree 2(3)+1=7. Wow!!

Let's develop the 2-point (m=1) Gaussian quadrature rule on [-1,1], by requiring that it has polynomial degree 2m+1=3.

$$G_2(f) = \omega_0 f(x_0) + \omega_1 f(x_1)$$

It seems reasonable to guess that by symmetry, $\aleph_0 = -\aleph_1$ and $\omega_0 = \omega_1$,

Then, requiring exact results for each element of the monomial basis for P_3, the constraints are:

$$G_2(1) = \omega_1 \cdot 1 + \omega_2 \cdot 1 = 2\omega_1 = \int_{-1}^{1} dx = 2 \Rightarrow \omega_1 = 1$$

So
$$G_2(f) = f(-x_1) + f(x_1)$$
.
 $G_2(x) = -x_1 + x_1 = 0$ need $\int x dx = 0$ for regardless $f(x)$, $f(x) = (-x_1)^2 + x_1^2 = 2x_1^2$ need $\int x^2 dx = \frac{2}{3} \Rightarrow x_1^2 = \frac{1}{3}$, $x_1 = \frac{1}{3}$.
Thus $G_2(f) = f(-\frac{1}{3}) + f(\frac{1}{3})$

And
$$G_{3}(x^{3}) = -x_{1}^{3} + x_{1}^{3} = 0 \stackrel{\text{red}}{=} \int_{-1}^{1} x^{3} dx = 0$$

Let's compare with the inferior 2-point rule called the "trapezoid rule":

$$T(f) = f(-1) + f(1)$$
.

Let's test the 2-point Gaussian quadrature rule we just derived

```
(to integrate a non-polynomial function)
from numpy import sqrt, exp
def G2(f):
    return f(-1/sqrt(3)) + f(1/sqrt(3))
def T(f):
    return f(-1) + f(1)
            # here is the function we want to integrate
myf = exp
e = exp(1)
exact = e - 1/e
qauss approx = G2(myf)
trap approx = T(myf)
gauss error = (gauss approx-exact)/exact
trap error = (trap approx -exact)/exact
print(trap_error)
0.31303528549933135
```

Gauss is 100 times better than Trapezoid in this example.

print(gauss_error)
-0.003278714921135362

How to obtain the Gauss rules for larger m?

Ingredients are a set of polynomials (p0, p1, p2, ...) with the following properties:

- (1) pn has degree n
- pn has degree n
 pi is orthogonal to pj if i does not equal j, with respect to the inner product $(P_i, P_j) = (P_i, P_j)$ pn has n distinct roots in [-1.1]
- (3) pn has n distinct roots in [-1,1]

$$(P_i, P_j) = \int_{-1}^{1} P_i P_j$$

We could develop such a set as follows (choose some basis and orthogonalize it using Gram-Schmidt):

Say, let
$$p_0(x) = 1$$

 $p_1(x) = x$
 $p_2(x) = x^2$ nope, not rethogonal to p_0
 $= x^2 - \frac{1}{3}$

etc., obtiaining the "Legendre polynomials".

A convenient representation of the Legendre polynomials (I claim) is

$$p_n(x) = \frac{1}{2^n!} (x^2 - 1)^{n(n)}$$

Does this satisfy properties (1),(2),(3)?

(1) pn is of degree n? Yes! It's a polynomial of degree 2n diff'd n times.

(2) Orthogonality?

$$\int_{-1}^{1} P_{i} P_{j}^{2} = \int_{-1}^{1} (x^{2}-1)^{i} (x^{2}-1)^{j} dx, \quad i \neq j \notin \text{wolog } i \neq j.$$
Integrating by ports:
$$\int_{-1}^{1} P_{i} P_{j}^{2} = (x^{2}-1)^{i} (i-1)(x^{2}-1)^{j} (j+1) - \int_{-1}^{1} (x^{2}-1)^{i} (i-1)(x^{2}-1)^{j} (j+1)$$

$$= O \text{ because } (x^{2}-1)^{i} \text{ has a multiple roof } \text{foot } \text{foot } \text{ot } x^{2}=1.$$

$$= -\int_{-1}^{1} (x^{2}-1)^{i} (i-1)(x^{2}-1)^{j} (j+1) dx$$

$$= \pm \int_{-1}^{1} (x^{2}-1)^{i} (x^{2}-1)^{j} (x^{2}-1)^{j} dx = 0$$

$$= \pm \int_{-1}^{1} (x^{2}-1)^{i} (x^{2}-1)^{j} (x^{2}-1)^{j} dx = 0$$

(3) n distinct roots in [-1,1]? ... Yes: Google it > Stack Exchange

we use the roots of p(m+1) as the nodes, and choose the weights to obtain degree m

we will in fact have a rule of degree 2m+1 - the Gauss(-Legendre) rule.

where
$$S$$
, R are polynomials, $deg S = m < m+1$, $deg R < m+1$

See example below

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Observation (1)

$$G(P) = G(R)$$
 Why?

$$G(P) = G(S p_{m+1}) + G(R)$$

Observation (2)

Observation (3)

Summarizing

$$G(\mathbf{P}) = G(\mathbf{R}) \stackrel{(2)}{=} \int \mathbf{R} = \int \mathbf{P}$$



Example:
$$m=1$$
, $m+1=2$, $2m+1=3$

$$P(x) = x^{5} + 4x^{2} + 3x + 5$$

$$P_{2}(x) = x^{2} - \frac{1}{3}$$

$$x^{2} - \frac{1}{3} x^{3} + 4x^{2} + 3x + 5$$

$$x^{3} - \frac{1}{3}x$$

$$x^{3}$$

Error bound for quadrature rules

Recall Theorem 4.6 p215:

If
$$f \in C^{m+1}[a,b]$$
 then $R(x) = f(x) - p(x) = f(x+1)!$ $f \in C^{m+1}[a,b]$ then $R(x) = f(x) - p(x) = f(x+1)!$ $f \in C^{m+1}[a,b]$ then $f \in C^{m+$

and we can reduce "H" by dividing our domain [a,b] into small subintervals or "panels", applying the rule to each panel and adding the results.

In practice, though, ... we estimate our error numerically, as follows ...

Suppose we have two rules, such as G7 and "K15",

where deg(K15) >> deg(G7).

Then
$$G_7(f) - K_{IS}(f)$$
 is a reasonable estimate of the error in G7(f).

K in K15 above stands for Kronrod.

The (2m+1)-point Kronrod rule is the (2m+1)-point rule of maximal polynomial degree that re-uses the sample points of the m-point Gaussian rule.

Let's test the 7,15-point Gauss-Kronrod pair of rules

```
from numpy import *
# nodes and weights for the interval [-1,1] copied-and-pasted from Wikipedia
knodes = array([
-0.9914553711208126392,
-0.9491079123427585245,
-0.8648644233597690728.
-0.7415311855993944399
-0.5860872354676911303.
-0.4058451513773971669.
-0.2077849550078984676,
0.2077849550078984676,
 0.4058451513773971669,
 0.5860872354676911303,
 0.7415311855993944399,
0.8648644233597690728,
 0.9491079123427585245,
0.9914553711208126392])
gnodes = knodes[1::2]
kweights = array([
0.022935322010529224964,
0.06309209262997855329.
0.10479001032225018384.
0.14065325971552591875,
0.16900472663926790283,
0.1903505780647854099,
0.20443294007529889241,
0.2094821410847278280,
0.20443294007529889241,
0.1903505780647854099,
0.16900472663926790283,
0.14065325971552591875
0.10479001032225018384,
0.06309209262997855329.
0.0229353220105292249641)
gweights = arrav([
0.12948496616886969327061143267908201832858740225995, \
0.2797053914892766679014677714237795824869250652266,
0.3818300505051189449503697754889751338783650835339,
0.4179591836734693877551020408163265306122448979592,
0.3818300505051189449503697754889751338783650835339,
0.2797053914892766679014677714237795824869250652266,
0.1294849661688696932706114326790820183285874022599 ])
```

```
def g7(f,a,b):
    global gnodes, gweights
    center = (a+b)/2.
    halfwidth = (b-a)/2.
    nodes = center + halfwidth*gnodes
   weights = halfwidth*gweights
    return dot( weights, f(nodes) )
def k15(f,a,b):
    global knodes,kweights
    center = (a+b)/2.
   halfwidth = (b-a)/2.
    nodes = center + halfwidth*knodes
   weights = halfwidth*kweights
    return dot( weights, f(nodes) )
# test them
def myf(x): return sin(x) #x**10
def myF(x): return -cos(x) # for this example we can provide an antiderivative to obtain the integral exactly
a,b = 2,5 # randomly chosen interval
exact = myF(b) - myF(a)
print( 'Exact
                  :',exact )
Gest = g7(myf,a,b)
print( "Gauss 7 : ", Gest, ', error:', abs(Gest-exact)/exact )
Kest = k15(myf,a,b)
print( "Kronrod 15:", Kest, ', error:', abs(Kest-exact)/exact )
```

```
Exact : -0.6998090220103687

Gauss 7 : -0.6998090220106843 , error: -4.510322044637135e-13

Kronrod 15: -0.6998090220103687 , error: -0.0
```

Adaptive quadrature

How to attain an integral approximation with error tolerance ${\mathcal E}$?

Apply first directly to [a,b] and obtain error estimate.

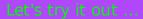
If error estimate $> \xi$,

then apply rule separately

on $\begin{bmatrix} a, \frac{a+b}{2} \end{bmatrix}$ with tolerance and on $\begin{bmatrix} a+b \\ 2 \end{bmatrix}$, b $\end{bmatrix}$ with tolerance $\begin{bmatrix} a+b \\ 2 \end{bmatrix}$

and add the results.

Do this recursively until total error estimate < \bigcirc .

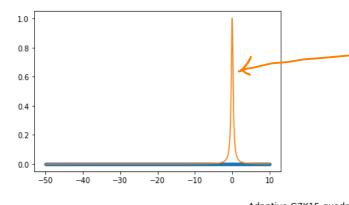




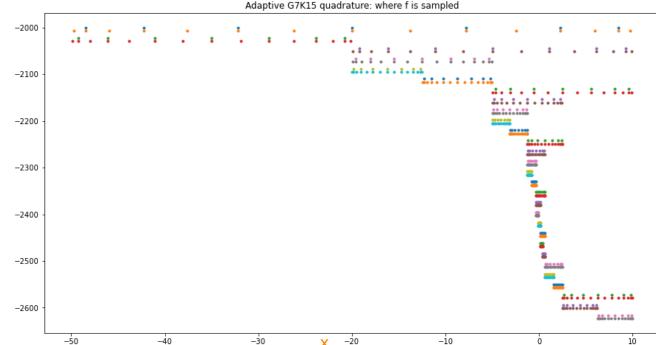
Adaptive quadrature with G7K15

```
import numpy as np
from qk import q7,k15
import matplotlib.pyplot as plt
count = 0
def runge(x):
    global count
    plt.plot(x,-count*np.ones\ like(x),'.') # plot the points where we've been asked to evaluate the function
    count += len(x)
    return 1/(1+12*x**2)
f = runge
def F(x): # antiderivative of Runge
    s = np.sqrt(12)
    return np.arctan(s*x)/s
def quad( f, a, b ):
    g = g7 (f, a, b)
    k = k15(f, a, b)
    error estimate = 200*abs(g-k)**1.5
    return k, error estimate
def adaptive( f, a, b, tol ):
    i,e = quad(f, a, b)
    if e <= tol:</pre>
        return i, e
    else:
        midpt = (a+b)/2
        i1, e1 = adaptive( f, a  , midpt, tol/2 )
        i2, e2 = adaptive( f, midpt, b , tol/2 )
        return i1+i2, e1+e2
a = -50
b = 10
exact = F(b) - F(a)
xx = linspace(a,b,2000); plt.figure(); plt.plot(xx,f(xx))
print( "exact ", exact )
plt.figure(figsize=(15,8))
i,e = adaptive( f, a, b, 1e-10 )
print( "approx", i, ", error estimate", e )
print( "actual error", i-exact )
plt.title('Adaptive G7K15 quadrature: where f is sampled')
plt.savefig('temp.pdf')
```





Integrand (Runge function) has a big spike in a small region of our interval [-50,10].



Closing question:

"How is it that the Gaussian rules are so remarkably accurate for non-polynomial functions?"

Consider the 2-point (m=1) Gaussian rule. There is a unique polynomial of degree 1 or less that interpolates the sampled values of the function f for which we want to evaluate \int_{-1}^{1} .

But there is an entire 2D family of functions within P_3 that interpolate the data, and if we choose the Gaussian sample points they all have the same integral (because we know the 2-point Gauss rule gets the integral exactly for all $p \in P_3$).

Thus there is an entire 2D family of interpolating cubics with the same integral, and it is plausible that some of them approximate f quite well, yielding an accurate approximation of its integral. Illustration below.

Let's take a look at an example ...