A final note on solving non-linear systems

Homotopy (morphing) in root-finding

Often hard to find a good enough starting point $x^{(\circ)}$ for Newton or quasi-Newton because the "basin of attraction" of the root is very small.

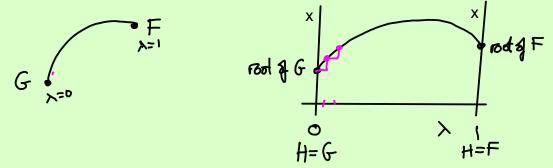
How to come up with a good × ?

A powerful idea is to consider a family of systems $H: \mathbb{R}^n \succeq [0, i] \to \mathbb{R}^n$ $H(x, \lambda)$

where H(x,0) = G(x), H(x,1) = F(x)

the function whose root we want some other "nicer, easier" function

Then what we do is start with an easily obtained root of G, and then gradually "morph" G into F, incrementing the parameter λ in $H(x,\lambda)$, tracking the root as we go, using the root on the previous value of λ as the starting x(0) for the current value of λ , until finally at $\lambda = 1$ we have a root of F.



Choosing a suitable homotopy is not necessarily easy.

A homotopy can easily fail, such as when the root of G and the sought root of F are not connected by a curve, maybe like this:

A lazy choice of homotopy, like G(x) = x - r (really easy to find a root (it's r!)) and $H(x,\lambda) = (1-x)G(x)$ is almost certain to fail.

Knowledge about the particular problem at hand will be useful in constructing a successful homotopy, such as in the toy suspension bridge system of Project Option 3. For example, a good approach might be to let G be a version of F with different, and particularly simplifying, parameter values, and let the homotopy be a linear ramp from the simplifying parameter values to the values you're actually interested in.

Ch. 9 Optimization

Given a scalar-valued function $+: D \subseteq \mathbb{R}^7 \to \mathbb{R}$

find a minimizer of f on D, that is find $\times \stackrel{*}{\in} D$ s.t. $f(\times \stackrel{*}{\times}) \leq f(\kappa) \ \forall \times \in D$ or find a local minimizer, that is find $x^* \in \mathbb{D}$ s.t. $f(x^*) \leq f(x) \forall x \in Some n' hood of <math>x^*$. If you want to maximize some function g, then define f = -g and minimize f.

Simplest case is n=1: $f: \mathbb{R} \rightarrow \mathbb{R}$ Ex:

Ex: finding the "best" Bezier approximation to a quarter-circle in Homework 6 Q7.



This problem is reminiscent of 1D root-finding, but a little trickier.

Recall the bisection method for root-finding, where we had a "bracket" [a,b] of the sought root: sign(f(a))!= sign(f(b)) guaranteeing a root in [a,b] if f is continuous, by IVT theorem. Idea was to shrink the bracket repeatedly until narrower than our error tolerance.

Can we do something similar for minimization?

Let's say that f is "unimodal" on [a,b] if $\exists c \in (a,b)$ such that

f is strictly increasing on [c,b].

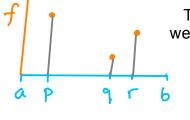
Then c is the unique minimizer of f on [a,b].

f is strictly decreasing on [a,c]

As analog of the "bracket" in root-finding, for a function f, let's define a "vee" as a triple of points (p,q,r)

such that p,q,r in [a,b] and

f(p) > f(q) and f(r) > f(q).



Then if f is unimodal on [a,b], we are guaranteed that

C ∈ (p, r).

(Why?)

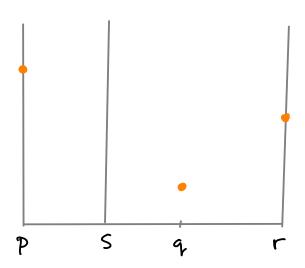
ni



The idea, analogously to bisection, is to repeatedly shrink the "vee", such that the width "|r-p|" of the vee goes to zero.

How to do that?

Let's pick a point $S \in (P, r)$, $S \neq \gamma$



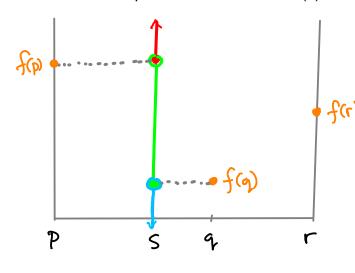
Your ideas on what to pick for s?

Idea 1: (p+r)/2

Idea 2: midpoint of one half, alternating

Idea 3: midpoint of larger subinterval

Regardless of the precise choice of s, suppose its between p and r, and consider the possible outcomes for f(s):



If f(s) ∈],

wait! Can't happen. Contradicts unimodality. If $f(s) \in \mathcal{J}$ a new smaller rea is (s, q, r)

a new smaller vee is (P, S, q).

If f(s)= = = f(q)

a new smaller vee is (5, t, q) where t is ybetween 58 g.

Let's find out what the guaranteed vee-width reduction is Idea 1: (p+r)/2 for your various ideas on choosing s.

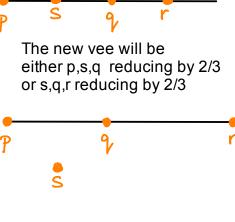
Idea 2: midpoint of one half, alternating

Idea 3: midpoint of larger subinterval



The new vee is either p,q,s reducing by 3/4 or q,s,r reducing by 1/2

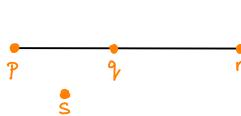




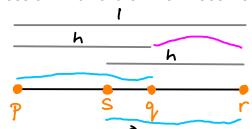
Best possible outcome would be to alternate 1/2 followed by a 2/3.

Average per-step reduction is $sqrt(1/2 . 2/3) = sqrt(1/3) \sim .577$ which is quite good - almost as good as bisection for root finding. Worst case is sqrt(2/3 . 3/4) = sqrt(1/2) < .7

best guarantee



Is there a choice for s where we are assured a specific reduction on each step? Maybe if we find a way to preserve the geometry, regardless of which of the 2 new vees we are forced to choose?



Is it possible to choose h such that the 2 possible vee-widths are the same?

We would need
$$\frac{h}{1} = \frac{(-h)}{h} \rightarrow h^2 + h - 1 = 0$$

$$h = \frac{-1 \pm \sqrt{1^2 - (-4)^2}}{2} = \frac{-1 + \sqrt{5}}{2} = \frac{\text{golden mean}}{2} = .618...$$

This is called "golden mean search" or "golden section search".

Start by setting q = (1-h)p + hr, that is q is a fraction h of the way from p to r.

With this scheme we are assured a width reduction of ~.618 at each step.

(Almost as good as bisection in root-finding.)

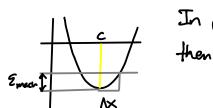
A caution about precision

We cannot demand nearly as high precision in optimization as in root-finding.

Recall in root-finding it's ok to set "tol" $\sim 10 \varepsilon_{\text{mech}} \sim 10^{-15}$.

 $\frac{a(x-c)^2+b-b}{b} = \frac{a(x-c)^2}{b}$ $\stackrel{\text{Set}}{=} \sum_{n=0}^{\infty} 1 : \Delta x = \sqrt{\frac{b}{a}} \sum_{n=0}^{\infty} n$

But for minimization, typically functions have quadratic minima, and near c, f changes hardly at all as x changes.



In fact, if
$$f(x) \approx \alpha(x-c)^2 + b$$
,
then if $f(x) - f(c) = \mathcal{E}_{mach}$, $\left| \frac{\Delta x}{c} \right| = \sqrt{\frac{b}{\alpha c^2}} \sqrt{\mathcal{E}_{mach}}$.
Or obsorbe $\left| \Delta x \right| = \sqrt{\frac{b}{a}} \sqrt{\mathcal{E}_{mach}}$.

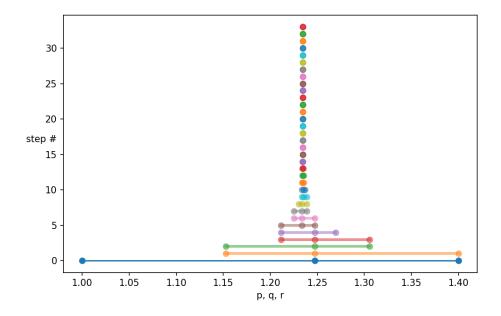
So our limit of precision in finding the minimum is

Don't ask for more!

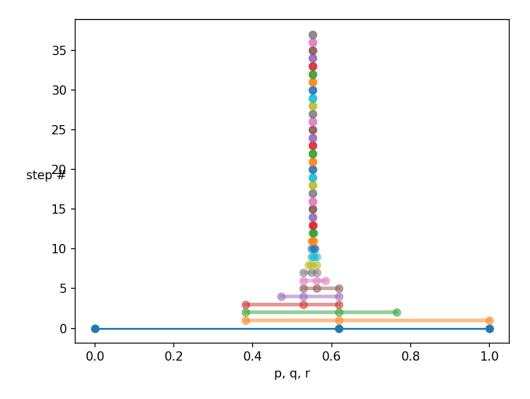
A guick implementation of golden section search ...

```
import numpy as np
import matplotlib.pyplot as plt
%matplotlib notebook
amp = 100
def golden(f,p,r,tol):
    h = (-1 + np.sqrt(5))/2
    q = (1-h)*p + h*r
    # verify that p,q,r is a vee
    assert( p<r and f(p)>f(q) and f(r)>f(q))
    count = 0
    plt.plot((p,q,r),count*np.ones(3),'o-')
    while r-p > tol:
        s = p + r - q
        fs, fq = f(s), f(q)
        if s<q:
            if fs<fq:
                q,r = s,q
            elif fs>fq:
                p = s
            else: #fs==fq
                p,r = s,q
                q = (1-h)*p + h*r
        else: #s>q:
            if fs<fq:
                p,q = q,s
            elif fs>fq:
                r = s
            else: #fs==fq
                p,r = q,s
                q = (1-h)*p + h*r
        count += 1
        plt.plot([p,q,r], count*np.ones(3), 'o-', color=f'C{count}', lw = 3, alpha=0.5)
def myf(x): return (x-1.23456789)**2 + 5 # an example function with a quadratic minimum
p,r = 1,1.4
x = np.linspace(p,r,200)
plt.figure(figsize=(8,5))
golden(myf, 1, 1.4, 1.e-8)
```

[space for results]



Now let's do Homework 6 Q7 (Bezier approx to quarter-circle) properly!



(0.5519149498173181, 0.5519149696417678)

Faster methods for smoother f

So far, we've depended only on f being unimodal. No smoothness (or even continuity) required.

We can do better (faster) if f has some smoothness.

For example, recall that Newton's method converges quadratically to a root of g if $g \in C^2$.

Now we are seeking a minimizer of f. If $f \in C^3$ then $g \equiv f' \in C^2$ and a local minimizer of f is a root of g to which Newton will converge quadratically:

$$x^{(k+1)} = x^{(k)} - \frac{f'(x^{(k)})}{f''(x^{(k)})}.$$

Another idea that doesn't depend on quite so much smoothness is **successive quadratic interpolation**.

Next class, we'll explore optimization with respect to a vector variable (n > 1), (which is a huge area).

